

## EXERCISE SHEET: PROPOSITIONAL PROOF THEORY

### Exercise 1: Natural Deduction

Prove the following formula using natural deduction.

$$\neg(\forall x(\exists y(\neg P(x) \wedge P(y))))$$

#### Solution

$$\frac{\frac{\frac{[\forall x(\exists y(\neg P(x) \wedge P(y)))]_1}{\exists y(\neg P(x_1) \wedge P(y))} \forall E \quad \frac{\frac{[\neg P(x_1) \wedge P(y_1)]_2}{P(y_1)} \wedge E_2 \quad \frac{\frac{[\forall x(\exists y(\neg P(x) \wedge P(y)))]_1}{\exists y(\neg P(y_1) \wedge P(y))} \forall E \quad \frac{[\neg P(y_1) \wedge P(y_2)]_3}{\neg P(y_1)} \wedge E_1}{\neg P(y_1)} \exists E(3)}{\perp} \neg E}{\neg(\forall x(\exists y(\neg P(x) \wedge P(y))))} \neg I(1)$$

### Exercise 2: Sequent Calculus

Prove the following formulae in sequent calculus:

1.  $\neg \exists x P(x) \rightarrow \forall x \neg P(x)$
2.  $(\forall x(P \vee Q(x))) \rightarrow (P \vee \forall x Q(x))$

#### Solution

$$\begin{array}{l} 1. \frac{\frac{\frac{\frac{P(y) \Rightarrow P(y), \exists x P(x)}{\Rightarrow P(y), \exists x P(x), \neg P(y)} \neg R}{\Rightarrow \exists x P(x), \neg P(y)} \exists R}{\Rightarrow \exists x P(x), \forall x \neg P(x)} \forall R}{\neg \exists x P(x) \Rightarrow \forall x \neg P(x)} \neg L}{\Rightarrow \neg \exists x P(x) \rightarrow \forall x \neg P(x)} \rightarrow R \\ \\ 2. \frac{\frac{\frac{\frac{\forall x(P \vee Q(x)), P \Rightarrow P, Q(x)}{\forall x(P \vee Q(x)), (P \vee Q(x)) \Rightarrow P, Q(x)} \forall L}{\forall x(P \vee Q(x)) \Rightarrow P, Q(x)} \forall L}{\frac{\frac{\frac{\forall x(P \vee Q(x)) \Rightarrow P, Q(x)}{\forall x(P \vee Q(x)) \Rightarrow P, \forall x Q(x)} \forall R}{\forall x(P \vee Q(x)) \Rightarrow P \vee \forall x Q(x)} \vee R}{\Rightarrow \forall x(P \vee Q(x)) \rightarrow P \vee \forall x Q(x)} \rightarrow R \end{array}$$

### Exercise 3: Natural Deduction can Simulate Sequent Calculus II

In exercise 6.2 we proved that if  $\Gamma \vdash_G \Delta$  then  $\Gamma \vdash_N \bigvee \Delta$  for formulae in propositional logic. Augment your proof by the new cases for FOL.

### Solution

**Case 1:**  $\frac{F[t/x], \forall x F, \Gamma \vdash_G \Delta}{\forall x F, \Gamma \Rightarrow \Delta} \forall L$

By IH we have  $T : F[t/x], \forall x F, \Gamma \vdash_N \bigvee \Delta$

We obtain our proof by replacing every occurrence of the open assumption  $F[t/x]$  in  $T$  by the proof tree  $\frac{\forall x F}{F[t/x]} \forall E$ .

**Case 2:**  $\frac{\Gamma \vdash_G F[y/x], \Delta}{\Gamma \Rightarrow \forall x F, \Delta} \forall R$  Where  $y \notin \text{fv}(\Gamma, \forall x F, \Delta)$

By IH we have  $T : \Gamma \vdash_N F[y/x] \vee \bigvee \Delta$

$$\frac{T \quad \frac{\frac{[F[y/x]]_1}{\forall x F} \forall I \quad \frac{[\bigvee \Delta]_1}{\forall x F \vee \bigvee \Delta} \forall I_2}{\forall x F \vee \bigvee \Delta} \forall I_1}{\forall x F \vee \bigvee \Delta} \forall E(1)$$

**Case 3:**  $\frac{F[y/x], \Gamma \vdash_G \Delta}{\exists x F, \Gamma \Rightarrow \Delta} \exists L$  where  $y \notin \text{fv}(F, \Gamma, \Delta)$

By IH we have  $T : F[y/x], \Gamma \vdash_N \bigvee \Delta$

$$\frac{\exists x F \quad T}{\bigvee \Delta} \exists E(1)$$

**Case 4:**  $\frac{\Gamma \vdash_G F[t/x], \exists x F, \Delta}{\Gamma \Rightarrow \exists x F, \Delta} \exists R$

By IH we have  $T : \Gamma \vdash_N F[t/x] \vee \exists x F \vee \bigvee \Delta$

$$\frac{T \quad \frac{\frac{[F[t/x]]_1}{\exists x F} \exists I \quad \frac{[\exists x F \vee \bigvee \Delta]_1}{\exists x F \vee \bigvee \Delta} \exists I_1}{\exists x F \vee \bigvee \Delta} \exists I}{\exists x F \vee \bigvee \Delta} \exists E(1)$$

### Exercise 4: Counterexamples from Sequent Calculus

Consider the statement  $\forall x(P(x) \rightarrow \neg P(f(x)))$ .

1. What happens when trying to prove the validity of this formula in sequent calculus?
2. How can we derive a countermodel from the proof tree?
3. Is there a smaller countermodel?

### Solution

The proof tree gets stuck:

$$\frac{\frac{\frac{P(y), P(f(y)) \Rightarrow}{P(y) \Rightarrow \neg P(f(y))} \neg R}{\Rightarrow P(y) \rightarrow \neg P(f(y))} \rightarrow R}{\Rightarrow \forall x (P(x) \rightarrow \neg P(f(x)))} \forall R$$

As in the lecture, we can create a countermodel  $\mathcal{A}$ : Let  $U_{\mathcal{A}}$  be the set of all terms over  $y, f(\cdot)$ , set  $y^{\mathcal{A}} := y$ ,  $f^{\mathcal{A}}(t) := f(t^{\mathcal{A}})$ , and  $P^{\mathcal{A}} := \{y, f(y)\}$ . Then  $\mathcal{A} \models P(y)$  and  $\mathcal{A} \models P(f(y))$  and hence  $\mathcal{A} \not\models \forall x (P(x) \rightarrow \neg P(f(x)))$ . Note that  $\mathcal{A}$  is infinite, but there are countermodels with just two elements  $\{a, b\}$ : Set  $f(a) := b$ ,  $f(b) := b$ ,  $P(a)$  and  $P(b)$ . Then  $P(a)$  and  $P(f(a)) = P(b)$ .