

## Petri nets — Exercise sheet 7

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### Exercise 7.1

- (a) Show that an S-system  $(N, M_0)$  is bounded for any  $M_0$ .
- (b) Show that if  $(N, M_0)$  is a live S-system and  $M'_0 \geq M_0$ , then  $(N, M'_0)$  is also live.
- (c) Give a 1-bounded S-system  $(N, M_0)$  where  $M_0(S) > 1$ .

### Exercise 7.2

- (a) Give a T-system that is bounded and connected but not strongly connected.
- (b) Give a strongly connected T-system  $(\mathcal{N}, M_0)$  which is not live and such that  $M_0 \neq \mathbf{0}$ .
- (c) Let  $(\mathcal{N}, M_0)$  be a T-system. Show that if  $(\mathcal{N}, M_0)$  is strongly connected and live, then it is bounded.
- (d) Reprove (c), but this time without assuming that  $(\mathcal{N}, M_0)$  is live.

### Exercise 7.3

- (a) Let  $\mathcal{N} = (P, T, F)$  be a Petri net, and let  $s, t \in T$  be such that  $\bullet s \cap t \bullet = \emptyset$ . Show that if  $M \xrightarrow{ts} M'$ , then  $M \xrightarrow{st} M'$ .
- (b) Let  $\mathcal{N} = (P, T, F)$  be a Petri net which is not strongly connected. Show that  $P \cup T$  can be partitioned into two disjoint sets  $U, V \subseteq P \cup T$  such that  $F \cap (V \times U) = \emptyset$ .
- (c) Let  $U$  and  $V$  be a partition as in (b). Show that if  $M \xrightarrow{\sigma} M'$ , then there exist  $\sigma_U \in (T \cap U)^*$  and  $\sigma_V \in (T \cap V)^*$  such that  $\sigma = \sigma_U \sigma_V$  and  $M \xrightarrow{\sigma_U \sigma_V} M'$ .
- (d) Let  $(\mathcal{N}, M_0)$  be live and bounded. Use (a), (b) and (c) to show that  $\mathcal{N}$  is strongly connected.

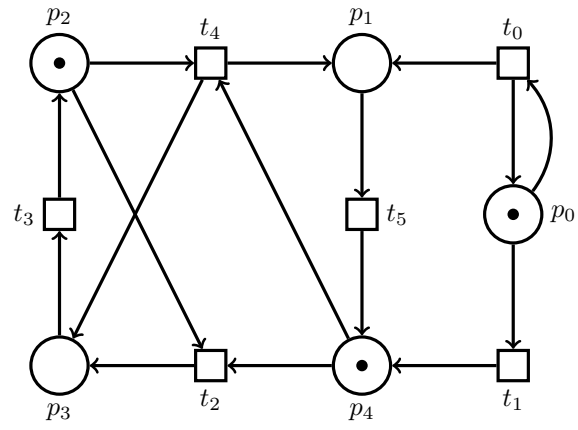
### Exercise 7.4

For each  $n \in \mathbb{N}$ , give a 1-bounded T-system  $(N, M_0)$  with  $n$  transitions and a reachable marking  $M$  such that the minimal occurrence sequence  $\sigma$  with  $M_0 \xrightarrow{\sigma} M$  has a length of  $\frac{n(n-1)}{2}$ .

*Hint:* First try find a Petri net and a marking for  $n = 3$ , where the minimal sequence has length 3. For this a net with 4 places suffices. Then try to generalize your solution.

### Exercise 7.5

Consider the following free-choice system  $(\mathcal{N}, M_0)$ :



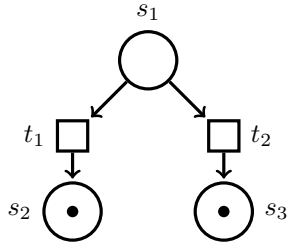
- (a) Give all minimal proper siphons of  $(\mathcal{N}, M_0)$ .
- (b) Use (a) to say whether  $(\mathcal{N}, M_0)$  is live or not.

**Exercise 7.6 (From [1, ex. 5.3])**

★ Exhibit a live and bounded free-choice system  $(\mathcal{N}, M_0)$  in which there exists a minimal trap which is not a siphon.

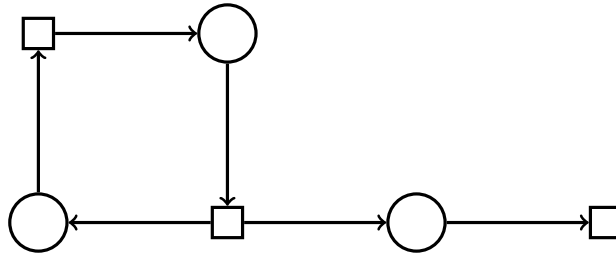
**Solution 7.1**

1. By the fundamental property of S-systems (Proposition 5.1.1), for every reachable marking, we have  $M(S) = M_0(S)$  and therefore  $M(s) \leq M_0(S)$  for all  $s \in S$ .
2. By the liveness theorem for S-systems (Theorem 5.1.3),  $(N, M_0)$  is live iff  $N$  is strongly connected and  $M_0(S) > 0$ , and as  $M'_0(S) \geq M_0(S) > 0$ ,  $(N, M'_0)$  is also live.
3. Due to the boundedness theorem for S-systems, the system needs to be non-live. The following Petri net is a non-live, 1-bounded S-system with  $M_0(S) > 1$ :

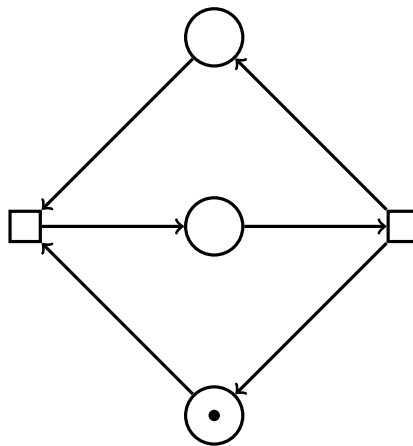


**Solution 7.2**

- (a) By Corollary 5.2.5, the system can not be live, and by Theorem 5.2.3. (Liveness Theorem for T-systems), it needs to have an unmarked circuit. The following system is then bounded, but not strongly connected:



- (b) By Theorem 5.2.3 (Liveness Theorem for T-systems), the system needs to have an unmarked circuit. The following system is then strongly connected, satisfies  $M_0 \neq \mathbf{0}$  and is not live:



- (c) Let  $\mathcal{N} = (P, T, F)$ . Let  $b = |M_0|$ . We show that every place is  $b$ -bounded. Let  $p \in P$ . Since  $\mathcal{N}$  is strongly connected,  $p$  lies on some circuit  $\gamma$ . Note that  $M_0(\gamma) \leq b$  and that  $(\mathcal{N}, M_0)$  is live. Therefore, by Theorem 5.2.4,  $p$  is  $b$ -bounded.  $\square$
- (d) Let  $\mathcal{N} = (P, T, F)$ . Let  $b = |M_0|$ . We show that every place is  $b$ -bounded. Let  $p \in P$ . Since  $\mathcal{N}$  is strongly connected,  $p$  lies on some circuit  $\gamma$ . By Proposition 5.2.2, for every reachable marking  $M$ ,  $M(\gamma) = M_0(\gamma)$ . So there can be no reachable marking  $M$  in which  $M(p) > b$  and  $p$  is  $b$ -bounded.

### Solution 7.3

- (a) Let  $X \in \mathbb{N}^P$  be such that  $M \xrightarrow{t} X \xrightarrow{s} M'$ . For the sake of contradiction, suppose  $s$  is not enabled in  $M$ . There exists  $p \in P$  such that  $p \in \bullet s$  and  $M(p) = 0$ . Since  $s$  is enabled in  $X$ , we have  $X(p) > 0$ . Therefore, it must be the case that  $p \in t^\bullet$ . This implies that  $p \in \bullet s \cap t^\bullet$  which is a contradiction. Thus,  $s$  is enabled in  $M$  and  $M \xrightarrow{s} Y$  for some marking  $Y \in \mathbb{N}^P$ .

Let us now show that  $t$  is enabled in  $Y$ . Let  $q \in \bullet t$ . We must show that  $Y(q) > 0$ .

Case 1:  $q \notin \bullet s$ . If  $q \notin \bullet s$ , then  $Y(q) \geq M(q) > 0$ .

Case 2:  $q \in \bullet s$ . If  $q \in \bullet s$ , then

$$Y(q) = M(q) - 1. \quad (1)$$

Since  $s$  is enabled in  $X$ , we have  $X(q) > 0$ . Moreover,  $q \notin t^\bullet$  since  $\bullet s \cap t^\bullet = \emptyset$ . This implies that  $M(q) > X(q)$ , and hence  $M(q) \geq 2$ . By (1), we derive  $Y(q) \geq 1$ .  $\square$

- (b) Since  $\mathcal{N}$  is not strongly connected, there exist  $u, v \in P \cup T$  such that there is no path from  $v$  to  $u$ . Let

$$\begin{aligned} U &= \{x \in P \cup T : \text{there is a path from } x \text{ to } u\}, \\ V &= (P \cup T) \setminus U. \end{aligned}$$

Note that both sets are non empty since  $u \in U$  and  $v \in V$ . Moreover,  $U \cap V = \emptyset$  and  $U \cup V = P \cup T$  by definition.

Let us show that  $F \cap (V \times U) = \emptyset$ . Assume there exists  $e \in F \cap (V \times U)$ . There exist  $x \in U$  and  $y \in V$  such that  $(y, x) \in F$ . Since  $x \in U$ , there exists a path  $\sigma$  from  $x$  to  $u$ . Therefore,  $(y, x)\sigma$  is a path from  $y$  to  $u$ . This implies that  $y \in U$  which is a contradiction.  $\square$

- (c) Let  $U' = T \cap U$  and  $V' = T \cap V$ . Let us first show that  $\bullet(U') \cap (V')^\bullet = \emptyset$ . For the sake of contradiction, assume there exist  $s \in V'$ ,  $t \in U'$  and  $q \in P$  such that  $q \in s^\bullet$  and  $q \in \bullet t$ . We have  $(s, q) \in F$  and  $(q, t) \in F$ . If  $q \in U$ , then by (b) and  $(s, q) \in F$ , we obtain a contradiction. Similarly, if  $q \in V$ , then  $(q, t) \in F$  yields a contradiction.

We now prove the claim by induction of  $|\sigma|$ . If  $|\sigma| = 0$ , it follows trivially. Assume that  $|\sigma| > 0$  and that the claim holds for firing sequences of length  $|\sigma| - 1$ . There exist  $\sigma' \in T^*$ ,  $s \in T$  and  $Y \in \mathbb{N}^P$  such that  $\sigma = \sigma's$  and

$$M \xrightarrow{\sigma'} X \xrightarrow{s} M'.$$

By induction hypothesis, there exists  $\pi_U \in (U')^*$  and  $\pi_V \in (V')^*$  such that  $M \xrightarrow{\pi_U \pi_V} X$ . If  $s \in V'$  or  $|\pi_V| = 0$ , then we are done. Otherwise, let  $\pi'_V \in (V')^*$  and  $t \in V'$  be such that  $\pi_V = \pi'_V t$ . Since  $\bullet(U') \cap (V')^\bullet = \emptyset$ , we can apply (a) and obtain

$$M \xrightarrow{\pi_U \pi'_V s} Y \xrightarrow{t} M'$$

for some  $Y \in \mathbb{N}^P$ . By induction hypothesis, there exist  $\gamma_U \in (U')^*$  and  $\gamma_V \in (V')^*$  such that

$$M \xrightarrow{\gamma_U \gamma_V} Y.$$

Let  $\sigma_U = \gamma_U$  and  $\sigma_V = \gamma_V t$ . We are done since  $\sigma_U \in (U')^*$ ,  $\sigma_V \in (V')^*$  and  $M \xrightarrow{\sigma_U \sigma_V} M'$ .  $\square$

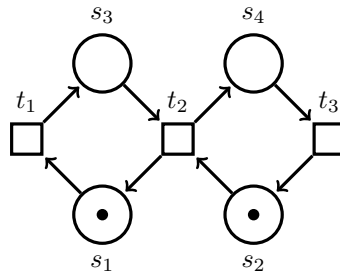
- (d) Let  $\mathcal{N} = (P, T, F)$ . For the sake of contradiction, assume  $\mathcal{N}$  is not strongly connected. By (b), there exists a partition  $U, V$  of  $P \cup T$  such that  $F \cap (V \times U) = \emptyset$ . Since  $\mathcal{N}$  is connected, there exist  $u \in U$  and  $v \in V$  such that  $(u, v) \in F$ . Let  $b \in \mathbb{N}$  be such that  $(\mathcal{N}, M_0)$  is  $b$ -bounded. Since  $(\mathcal{N}, M_0)$  is live, there exist  $\sigma \in T^*$  and  $M \in \mathbb{N}^P$  such that  $M_0 \xrightarrow{\sigma} M$  and  $(u, v)$  is taken  $b+1$  times. By (c), there exist  $\sigma_U \in U^*$  and  $\sigma_V \in V^*$  such that  $M_0 \xrightarrow{\sigma_U \sigma_V} M$ . Let  $X \in \mathbb{N}^P$  be such that  $M_0 \xrightarrow{\sigma_U} X \xrightarrow{\sigma_V} M$ .

Case 1:  $u \in P, v \in T$ . Since  $F \cap (V \times U) = \emptyset$ , there is no transition of  $V$  that puts tokens into places of  $U$ . Note that  $v$  decreases the amount of token of  $u$  by 1. Since  $X \xrightarrow{\sigma_V} M$ , these two observations imply that  $X(u) \geq b+1$ . As  $X$  is reachable from  $M_0$ , this contradicts  $(\mathcal{N}, M_0)$  being  $b$ -bounded.

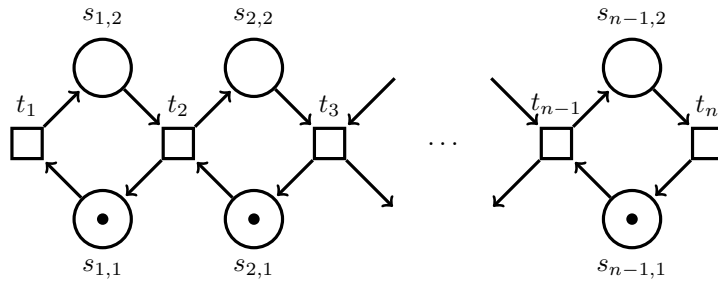
Case 2:  $u \in T, v \in P$ . Since  $F \cap (V \times U) = \emptyset$ , there is no transition of  $U$  that consumes tokens from places of  $V$ . Note that  $u$  increases the amount of token of  $u$  by 1. Since  $M_0 \xrightarrow{\sigma_U} X$ , these two observations imply that  $X(u) \geq b+1$ . This contradicts  $(\mathcal{N}, M_0)$  being  $b$ -bounded.  $\square$

**Solution 7.4**

For  $n = 3$ , we can take the following net with the marking  $M = (0, 0, 1, 1)$ . To reach this marking, we need to fire  $t_1$  and  $t_2$  to mark  $s_3$  and  $s_4$ . However, firing  $t_2$  undoes the effect of  $t_1$  on  $s_3$ , so we need to fire  $t_1$  twice. The minimal sequence is then  $\sigma = t_1 t_2 t_1$  of length 3.



This construction can be repeated for arbitrary  $n$ , as shown in the following sketch of a Petri net. To reach the marking  $M$  with  $M(s_{i,1}) = 0$  and  $M(s_{i,2}) = 1$  for all  $1 \leq i \leq n - 1$  with a minimal sequence, we need to fire  $\sigma = t_1 t_2 \dots t_{n-1} t_1 t_2 \dots t_{n-2} \dots t_1$ , which has a length of  $\sum_{i=1}^{n-1} i = \frac{n(n-1)}{2}$ .



**Solution 7.5**

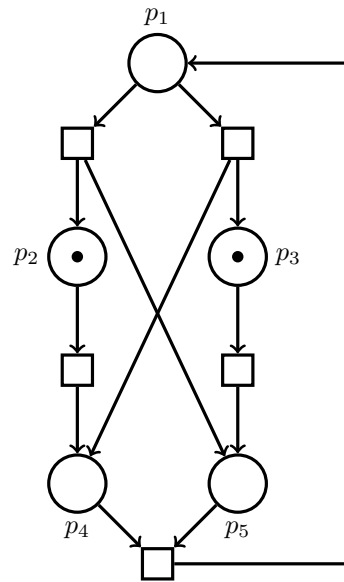
- (a) We claim that the system has two minimal proper siphons:  $\{p_0\}$  and  $\{p_2, p_3\}$ .

Let us show the claim. By inspecting  $\bullet p$  and  $p \bullet$  for every place  $p$ , we find a single siphon of size one:  $\{p_0\}$ . Moreover, we have  $\bullet\{p_2, p_3\} = \{t_2, t_3, t_4\} = \{p_2, p_3\} \bullet$ . Now, note that  $t_0 \in \bullet p_1$  and  $\bullet t_0 = \{p_0\}$ . Therefore, any siphon containing  $p_1$  must also contain  $p_0$ . Similarly, any siphon containing  $p_4$  must also contain  $p_0$ . Thus, no minimal siphon contains  $p_1$  or  $p_4$ , and we are done.  $\square$

- (b) The system is not live. By Commoner's Theorem, the system is live if and only if every minimal proper siphon contains a trap marked at  $M_0$ . The minimal siphon  $\{p_2, p_3\}$  is also a trap and it is marked at  $M_0$ . However, the minimal siphon  $\{p_0\}$  is not a trap and hence it does not contain a marked trap.

**Solution 7.6 (From [1, ex. 5.3])**

The following is a live and bounded free-choice system. The set  $\{p_1, p_4, p_5\}$  is a minimal trap, but not a siphon.



## References

- [1] Jörg Desel and Javier Esparza. *Free Choice Petri Nets*. Cambridge University Press, USA, 1995.