

# Commoner's Theorem

← free-choice net

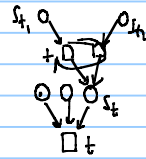
Theorem If every siphon of  $N$  contains a trap marked at  $M_0$ , then  $(N, M_0)$  is live

Proof Let  $M$  be a marking of  $N$ . A transition is

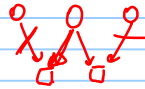
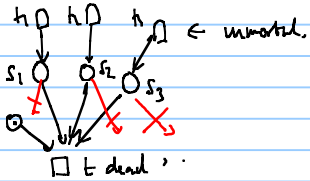
- **dead at  $M$**  if it is not enabled at any marking reachable from  $M$ .
- **live at  $M$**  if it "cannot die", i.e. it is not **(immortal)** dead at any marking reachable from  $M$ .
- **mortal at  $M$**  if there is  $M \xrightarrow{*} M'$  such that  $t$  is dead at  $M'$ .

• Claim: there is a reachable marking  $M$  such that every transition is either dead or immortal at  $M$ .

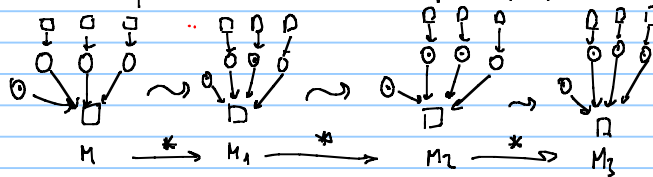
- Claim: for every  $t$  dead at  $M$  there exists  $S_t \in {}^*t$  s.t.
- $M(S_t) = 0$  and
  - every  $t' \in {}^*S_t$  is dead at  $M$ .



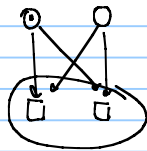
Proof of the claim: Let  $t$  be dead at  $M$  and let  $\{s_1, \dots, s_n\}$  be the places in  ${}^*t$  not marked at  $M$ . Assume that for every  $s_i$  there is  $t_i \in {}^*s_i$  live at  $M$ .



Since  $N$  is free-choice there are markings  $M_1, \dots, M_n$  s.t.



Contradiction that  $t$  is dead at  $M$ .



Theorem If  $(N, M_0)$  is free-choice and live, then every siphon of  $N$  contains a trap marked at  $M_0$ .

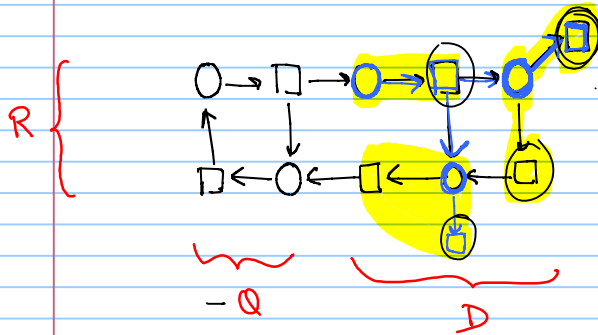
Proof sketch Let  $R$  be a siphon, and let  $Q$  be the maximal trap included in  $R$ . Let  $D = R \setminus Q$

Assume  $Q$  is initially unmarked.

We find a firing sequence  $\sigma$  that "empties"  $D$  without adding tokens to  $Q$ . After the execution of  $\sigma$  " $R$  is empty" which implies that  $(N, M_0)$  is not live.

The sequence  $\sigma$  is constructed as follows:

- we construct an allocation that assigns to each place  $S \in D$  a transition of  $S^*$ . The firing sequence  $\sigma$  only executes allocated transitions.



We need to guarantee:

- the allocation does not define cycles
- the allocation does not allocate any transition of  $Q$  (this would mark the trap)
- while there are tokens in  $D$  we can always fire allocated transitions again and again

## Cluster

Definition Let  $x$  be a place or transition of a net  $N = (S, T, F)$ . The cluster of  $x$ , denoted by  $[x]$  is the minimal set of nodes such that

- $x \in [x]$
- if  $s \in [x] \cap S$  then  $s^* \in [x]$
- if  $t \in [x] \cap T$  then  ${}^*t \in [x]$

Proposition Every node of a net belongs to exactly one cluster

→ The set of clusters is a partition of  $S \cup T$

## Definition

Let  $N = (S, T, F)$  be a net, and let  $C$  be a set of clusters of  $N$ . An **allocation** of  $C$  is a function  $\alpha: C \rightarrow T$  satisfying  $\alpha(c) \in c$

Proposition Let  $N$  be a free-choice net, let  $R$  be a set of places. Let  $Q$  be the maximal trap included in  $R$ , and let  $D = Q \cap R$ .

Let  $C = \{ [t] \mid t \in D^* \}$

There exists a circuit-free allocation

$\alpha: C \rightarrow T$  such that  $\alpha(C) \cap Q = \emptyset$

(circuit-free: the set of arcs  $\{ (s, \alpha([s])) \mid s \in D \} \cup F \cap (T \times S)$  does not contain any circuit)

Proof By induction on  $|R|$ .

•  $|R| = 0$ . Then  $C = \emptyset$  ✓

•  $|R| > 0$ . If  $R$  is a trap then  $C = \emptyset$  <sup>"way-out"</sup>

If  $R$  is not a trap: there is  $t \in R^* \setminus R$

Let  $R' = R \setminus {}^*t$ ,  $Q'$  be the maximal trap of  $R'$ ,

$D' = R' \cap Q'$ ,  $C' = \{ [t] \mid t \in D'^* \}$

By induction hypothesis there exists  $\alpha': C' \rightarrow T$  circuit-free for  $D'$  such that  $\alpha'(C') \cap Q' = \emptyset$ .

Define  $\alpha: C \rightarrow T$  by

$$\alpha(c) = \begin{cases} \alpha'(c) & \text{if } c \neq [t] \\ t & \text{if } c = [t] \end{cases}$$

We have to prove:

-  $\alpha(C) \subseteq \alpha'(C') \cup \{t\}$  ( $\alpha$  is well defined)

-  $\alpha$  is circuit-free for  $D$

-  $\alpha$  puts no tokens in  $Q$ .

Proposition Let  $\alpha$  be an allocation of  $a$

free-choice system

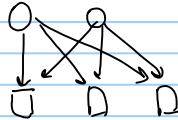
with domain  $C$ . There exists an infinite occurrence sequence  $\sigma$  that

- fires allocated transitions infinitely often

AND

- never fires any non-allocated transition of  $C$

Proof Immediate consequence of



Corollary The <sup>non</sup>reachability problem for free-choice nets is NP-complete

Proof

• NP-hardness: by reduction from SAT

• Membership in NP: - guess a path (or set it as advice)

- compute the largest trap in the path

- check the trap is empty at

$M_0$